# Project 1

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https://github.uio.no/FYS3150-G2-2023/Project-1

## PROBLEM 1

First, we rearrange the equation.

$$-\frac{d^2u}{dx^2} = 100e^{-10x}$$
$$\frac{d^2u}{dx^2} = -100e^{-10x}$$

Now we find u(x).

$$u(x) = \int \int \frac{d^2u}{dx^2} dx^2$$

$$= \int \int -100e^{-10x} dx^2$$

$$= \int \frac{-100e^{-10x}}{-10} + c_1 dx$$

$$= \int 10e^{-10x} + c_1 dx$$

$$= \frac{10e^{-10x}}{-10} + c_1 x + c_2$$

$$= -e^{-10x} + c_1 x + c_2$$

Using the boundary conditions, we can find  $c_1$  and  $c_2$ 

$$u(0) = 0$$

$$-e^{-10.0} + c_1 \cdot 0 + c_2 = 0$$

$$-1 + c_2 = 0$$

$$c_2 = 1$$

$$u(1) = 0$$

$$-e^{-10 \cdot 1} + c_1 \cdot 1 + c_2 = 0$$

$$-e^{-10} + c_1 + c_2 = 0$$

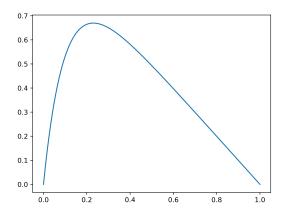
$$c_1 = e^{-10} - c_2$$

$$c_1 = e^{-10} - 1$$

Using the values that we found for  $c_1$  and  $c_2$ , we get

$$u(x) = -e^{-10x} + (e^{-10} - 1)x + 1$$
$$= 1 - (1 - e^{-10}) - e^{-10x}$$

The code for generating the points and plotting them can be found under. Point generator code: https://github.uio.no/FYS3150-G2-203/Project-1/blob/main/src/analyticPlot.cpp Plotting code: https://github.uio.no/FYS3150-G2-2023/Project-1/blob/main/src/analyticPlot.py Here is the plot of the analytical solution for u(x).



## PROBLEM 3

To derive the discretized version of the Poisson equation, we first need the Taylor expansion for u(x) around x for x + h and x - h.

$$u(x+h) = u(x) + u'(x)h + \frac{1}{2}u''(x)h^2 + \frac{1}{6}u'''(x)h^3 + \mathcal{O}(h^4)$$

$$u(x - h) = u(x) - u'(x)h + \frac{1}{2}u''(x)h^2 - \frac{1}{6}u'''(x)h^3 + \mathcal{O}(h^4)$$

If we add the equations above, we get this new equation:

$$u(x+h) + u(x-h) = 2u(x) + u''(x)h^{2} + \mathcal{O}(h^{4})$$

$$u(x+h) - 2u(x) + u(x-h) + \mathcal{O}(h^{4}) = u''(x)h^{2}$$

$$u''(x) = \frac{u(x+h) - 2u(x) + u(x-h)}{h^{2}} + \mathcal{O}(h^{2})$$

$$u''_{i}(x) = \frac{u_{i+1} - 2u_{i} + u_{i-1}}{h^{2}} + \mathcal{O}(h^{2})$$

We can then replace  $\frac{d^2u}{dx^2}$  with the RHS (right-hand side) of the equation:

$$-\frac{d^2u}{dx^2} = f(x)$$
$$\frac{-u_{i+1} + 2u_i - u_{i-1}}{h^2} + \mathcal{O}(h^2) = f_i$$

And lastly, we leave out  $\mathcal{O}(h^2)$  and change  $u_i$  to  $v_i$  to differentiate between the exact solution and the approximate solution, and get the discretized version of the equation:

$$\frac{-v_{i+1} + 2v_i - v_{i-1}}{h^2} = 100e^{-10x_i}$$

The value of  $u(x_0)$  and  $u(x_1)$  is known, using the discretized equation we solve for  $\vec{v}$ . This will result in a set of equations

$$-v_0 + 2v_1 - v_2 = h^2 \cdot f_1$$

$$-v_1 + 2v_2 - v_3 = h^2 \cdot f_2$$

$$\vdots$$

$$-v_{m-2} + 2v_{m-1} - v_m = h^2 \cdot f_{m-1}$$

where  $v_i = v(x_i)$  and  $f_i = f(x_i)$ . Rearranging the first and last equation, moving terms of known boundary values to the RHS

$$2v_1 - v_2 = h^2 \cdot f_1 + v_0$$
$$-v_1 + 2v_2 - v_3 = h^2 \cdot f_2$$
$$\vdots$$
$$-v_{m-2} + 2v_{m-1} = h^2 \cdot f_{m-1} + v_m$$

We now have a number of linear eqations, corresponding to the number of unknown values, which can be represented as an augmented matrix

$$\begin{bmatrix} 2v_1 & -v_2 & 0 & \dots & 0 & g_1 \\ -v_1 & 2v_2 & -v_3 & 0 & & g_2 \\ 0 & -v_2 & 2v_3 & -v_4 & & g_2 \\ \vdots & & \ddots & \vdots & \vdots \\ 0 & & -v_{m-2} & 2v_{m-1} & g_{m-1} \end{bmatrix}$$

Since the boundary values are equal to 0 the RHS can be renamed  $g_i = h^2 f_i$  for all i. An augmented matrix can be represented as  $\mathbf{A}\vec{x} = \vec{b}$ . In this case  $\mathbf{A}$  is the coefficient matrix with a tridiagonal signature (-1, 2, -1) and dimension  $n \times n$ , where n = m - 2.

**a**)

b)

#### PROBLEM 6

**a**)

Renaming the sub-, main-, and supdiagonal of matrix  $\boldsymbol{A}$ 

$$\vec{a} = [a_2, a_3, ..., a_{n-1}, a_n]$$
  
$$\vec{b} = [b_1, b_2, b_3, ..., b_{n-1}, b_n]$$
  
$$\vec{c} = [c_1, c_2, c_3, ..., c_{n-1}]$$

Following Thomas algorithm for gaussian elimination, we first perform a forward sweep followed by a backward sweep to obtain  $\vec{v}$ 

## Algorithm 1 General algorithm

```
procedure FORWARD SWEEP(\vec{a}, \vec{b}, \vec{c})
       n \leftarrow \text{length of } \vec{b}
      \vec{\hat{b}}, \, \vec{\hat{g}} \leftarrow \text{vectors of length } n.
      \hat{b}_1 \leftarrow b_1
                                                                                                                                                    ▶ Handle first element in main diagonal outside loop
      \hat{g}_1 \leftarrow g_1 for i = 2, 3, ..., n do
             d \leftarrow \frac{a_i}{\hat{b}_{i-1}}
                                                                                                                                                                                            ▷ Calculating common expression
      \hat{b}_i \leftarrow b_i - d \cdot c_{i-1}
\hat{g}_i \leftarrow g_i - d \cdot \hat{g}_{i-1}
\mathbf{return} \ \hat{b}, \ \hat{g}
procedure Backward sweep(\hat{b}, \hat{g})
      n \leftarrow \text{length of } \vec{\hat{b}}
       \vec{v} \leftarrow \text{vector of length } n.
      v_n \leftarrow \frac{\hat{g}_n}{\hat{b}_n}
      for i = n - 1, n - 2, ..., 1 do
             v_i \leftarrow \frac{\hat{g}_i - c_i \cdot v_{i+1}}{\hat{c}}
      return \vec{v}
```

b)

Counting the number of FLOPs for the general algorithm by looking at one procedure at a time. For every iteration of i in forward sweep we have 1 division, 2 multiplications, and 2 subtractions, resulting in 5(n-1) FLOPs. For backward sweep we have 1 division, and for every iteration of i we have 1 subtraction, 1 multiplication, and 1 division, resulting in 3(n-1) + 1 FLOPs. Total FLOPs for the general algorithm is 8(n-1) + 1.

## PROBLEM 8

## PROBLEM 9

**a**)

The special algorithm does not require the values of all  $a_i$ ,  $b_i$ ,  $c_i$ . We find the values of  $\hat{b}_i$  from simplifying the general case

$$\hat{b}_{i} = b_{i} - \frac{a_{i} \cdot c_{i-1}}{\hat{b}_{i-1}}$$

$$\hat{b}_{i} = 2 - \frac{1}{\hat{b}_{i-1}}$$

Calculating the first values to see a pattern

$$\hat{b}_1 = 2$$

$$\hat{b}_2 = 2 - \frac{1}{2} = \frac{3}{2}$$

$$\hat{b}_3 = 2 - \frac{1}{\frac{3}{2}} = \frac{4}{3}$$

$$\hat{b}_4 = 2 - \frac{1}{\frac{4}{3}} = \frac{5}{4}$$

$$\vdots$$

$$\hat{b}_i = \frac{i+1}{i}$$

for i = 1, 2, ..., n

## Algorithm 2 Special algorithm

 $\begin{aligned} & \text{procedure Forward sweep}(\vec{b}) \\ & n \leftarrow \text{length of } \vec{b} \\ & \hat{\vec{b}}, \ \vec{\hat{g}} \leftarrow \text{vectors of length } n. \\ & \hat{b}_1 \leftarrow 2 \\ & \hat{g}_1 \leftarrow g_1 \\ & \text{for } i = 2, 3, ..., n \text{ do} \\ & \hat{b}_i \leftarrow \frac{i+1}{i} \\ & \hat{g}_i \leftarrow g_i + \frac{\hat{g}_{i-1}}{\hat{b}_{i-1}} \\ & \text{return } \vec{b}, \ \vec{\hat{g}} \end{aligned}$   $& \text{procedure Backward sweep}(\vec{b}, \ \vec{\hat{g}})$   $& n \leftarrow \text{length of } \vec{\hat{b}} \\ & \vec{v} \leftarrow \text{vector of length } n. \\ & v_n \leftarrow \frac{\hat{g}_n}{\hat{b}_n} \\ & \text{for } i = n-1, n-2, ..., 1 \text{ do} \\ & v_i \leftarrow \frac{\hat{g}_i + v_{i+1}}{\hat{b}_i} \\ & \text{return } \vec{v} \end{aligned}$ 

▶ Handle first element in main diagonal outside loop

b)

For every iteration of i in forward sweep we have 2 divisions, and 2 additions, resulting in 4(n-1) FLOPs. For backward sweep we have 1 division, and for every iteration of i we have 1 addition, and 1 division, resulting in 2(n-1) + 1 FLOPs. Total FLOPs for the special algorithm is 6(n-1) + 1.